

Parallel robust solution of triangular equations

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What is a robust algorithm?

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A robust algorithm returns a result which

is free of **Inf** and **NaN**

and can be evaluated in the context of the user's application.

Example 1: $\mathbf{Ax} = \mathbf{b}$

$$\mathbf{Ax} := \begin{bmatrix} 1 & & & & \\ -1 & 1 & & & \\ -1 & -1 & 1 & & \\ -1 & -1 & -1 & 1 & \\ -1 & -1 & -1 & -1 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 2 \\ 4 \\ 8 \\ 16 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix} =: \mathbf{b}$$

Example 2: $SX = X\Lambda$

Matrix:

$$S = \begin{bmatrix} 1 & & & & \\ -5 & 2 & & & \\ -5 & -5 & 3 & & \\ -5 & -5 & -5 & 4 & \\ -5 & -5 & -5 & -5 & 5 \end{bmatrix}$$

Eigenvectors:

$$X = \begin{bmatrix} 1 & & & & \\ 5 & 1 & & & \\ 15 & 5 & 1 & & \\ 35 & 15 & 5 & 1 & \\ 70 & 35 & 15 & 5 & 1 \end{bmatrix}$$

Example 3: $\mathbf{AX} + \mathbf{XA}^T = \mathbf{bb}^T$

Matrices:

$$\mathbf{A} = \begin{bmatrix} \frac{1}{2} & & & & \\ -1 & \frac{1}{2} & & & \\ -1 & -1 & \frac{1}{2} & & \\ -1 & -1 & -1 & \frac{1}{2} & \\ -1 & -1 & -1 & -1 & \frac{1}{2} \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

Solution:

$$\mathbf{X} = \begin{bmatrix} 1 & 2 & 4 & 8 & 16 \\ 2 & 5 & 12 & 28 & 64 \\ 4 & 12 & 33 & 86 & 216 \\ 8 & 28 & 86 & 245 & 664 \\ 16 & 64 & 216 & 664 & 1921 \end{bmatrix}$$

Robust triangular solvers in LAPACK

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Subroutine	Problem	Comment
<code>xlatrs</code>	$\mathbf{Ax} = \alpha \mathbf{b}$	written 1991
<code>xtrevc</code>	$\mathbf{AX} = \mathbf{X}\Lambda$	
<code>xtgevc</code>	$\mathbf{AX} = \mathbf{BX}\Lambda$	
<code>xtrsyl</code>	$\mathbf{AX} + \mathbf{XB} = \alpha \mathbf{C}$	

Robust triangular solvers in ScaLAPACK?

From the source code of `pztrevc`¹, lines 186:190

The algorithm used in this program is basically backward (forward) substitution. It is the hope that scaling would be used to make the the code robust against possible overflow. But scaling has not yet been implemented in PZLATTRS which is called by this routine to solve the triangular systems. PZLATTRS just calls PZTRSV.

¹http://www.netlib.org/scalapack/explore-html/d5/dbc/pztrevc_8f_source.html Retrieved: October 19th 2019.

Robust backward substitution, $\mathbf{T}\mathbf{x} = \alpha\mathbf{b}$

Algorithm 1: $[\alpha, \mathbf{x}] = \text{RobustBackSub}(\mathbf{T}, \mathbf{b})$

```
1  $\alpha \leftarrow 1$ ;  
2 for  $j \leftarrow n, n-1, \dots, 1$  do  
3    $\beta = \text{ProtectDivision}(x_j, t_{jj})$ ;  
4   if  $\beta \neq 1$  then  
5      $\mathbf{x} \leftarrow \beta\mathbf{x}, \alpha \leftarrow \beta\alpha$ ;  
6    $x_j \leftarrow x_j/t_{jj}$ ;  
7   if  $j > 1$  then  
8      $\beta = \text{ProtectUpdate}(x_{\max}, \|\mathbf{T}_{1:j-1,j}\|_{\infty}, |x_j|)$ ;  
9     if  $\beta \neq 1$  then  
10       $\mathbf{x} \leftarrow \beta\mathbf{x}, \alpha \leftarrow \beta\alpha$ ;  
11       $\mathbf{x}_{1:j-1} \leftarrow \mathbf{x}_{1:j-1} - \mathbf{T}_{1:j-1,j}x_j$ ;  
12       $x_{\max} \leftarrow \|\mathbf{x}_{1:j-1}\|_{\infty}$ ;  
13 return  $[\alpha, \mathbf{x}]$ ;
```

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$$|b| \leq |t|\Omega.$$

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ProtectDivision

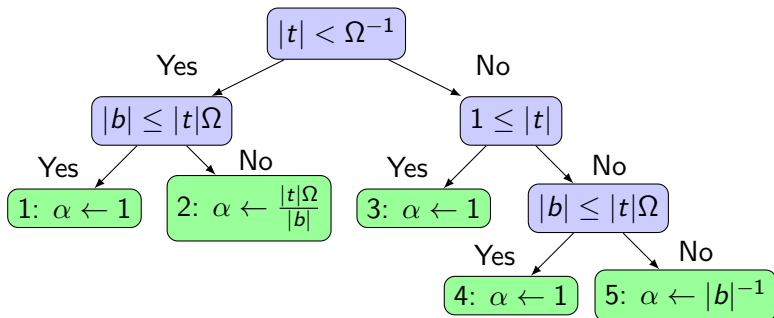
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Solution:



ProtectUpdate

The linear update $\mathbf{y} \leftarrow \mathbf{y} - \mathbf{T}\mathbf{x}$ is safe if

$$\|\mathbf{y}\|_{\infty} + \|\mathbf{T}\|_{\infty}\|\mathbf{x}\|_{\infty} \leq \Omega.$$

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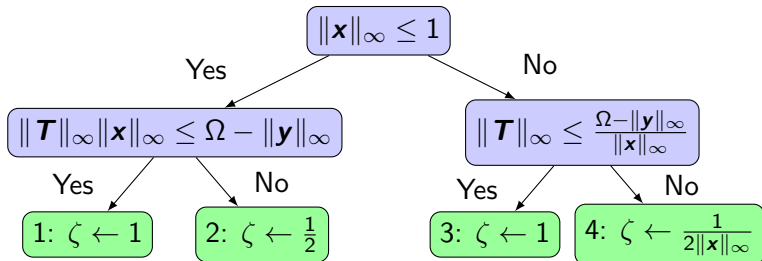
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How to develop blocked, parallel and robust codes?

The main idea:

Different processing units can use different units of measurement ...

Step 1: Partition the problem

$$\begin{bmatrix} T_{11} & T_{12} & \dots & T_{1M} \\ & T_{22} & \dots & T_{2m} \\ & & \ddots & \vdots \\ & & & T_{MM} \end{bmatrix} \begin{bmatrix} X_{11} & X_{12} & \dots & X_{1N} \\ X_{21} & \ddots & \ddots & X_{2N} \\ \vdots & \ddots & \ddots & \vdots \\ X_{M1} & \dots & \dots & X_{MN} \end{bmatrix} = \begin{bmatrix} B_{11} & B_{12} & \dots & B_{1N} \\ B_{21} & \ddots & \ddots & B_{2N} \\ \vdots & \ddots & \ddots & \vdots \\ B_{M1} & \dots & \dots & B_{MN} \end{bmatrix}$$

Step 2: Augment each block of data

A separate scaling factor α_{ijk} for each column of every block \mathbf{X}_{ij} .

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An augmented block

$$\langle \boldsymbol{\alpha}, \mathbf{X} \rangle$$

where

$$\boldsymbol{\alpha} = (\alpha_1, \alpha_2, \dots, \alpha_n), \quad \alpha_j \in (0, 1], \quad \mathbf{X} \in \mathbb{R}^{m \times n}$$

is used to represents the matrix

$$\mathbf{Y} \in \mathbb{R}^{m \times n}$$

where

$$\mathbf{y}_j = \alpha_j^{-1} \mathbf{x}_j, \quad j = 1, 2, \dots, n.$$

Step 3: Replace kernels with robust kernels ...

Problem: Given $\langle \alpha, \mathbf{X} \rangle$ and $\langle \beta, \mathbf{Y} \rangle$ find $\langle \nu, \mathbf{Z} \rangle$ such that

$$\nu_j^{-1} \mathbf{z}_j = \beta_j^{-1} \mathbf{y}_j - \alpha_j^{-1} \mathbf{T} \mathbf{x}_j, \quad j = 1, 2, \dots, n$$

using BLAS level-3 operations.

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Algorithm 3: $\langle \nu, \mathbf{Z} \rangle \leftarrow \text{RobustGEMM}(\langle \beta, \mathbf{Y} \rangle, \mathbf{T}, \langle \alpha, \mathbf{X} \rangle)$

- 1 **for** $j \leftarrow 1, 2, \dots, n$ **do**
 - 2 $\gamma_j \leftarrow \min\{\alpha_j, \beta_j\};$
 - 3 $\zeta_j \leftarrow \text{ProtectUpdate}((\gamma_j/\beta_j)\|\mathbf{y}_j\|_\infty, \|\mathbf{T}\|_\infty, (\gamma_j/\alpha_j)\|\mathbf{x}_j\|_\infty);$
 - 4 $\nu_j \leftarrow \zeta_j \gamma_j;$
 - 5 $\mathbf{Z} \leftarrow [\mathbf{Y} \text{diag}(\nu \circ \beta)] - \mathbf{T} [\mathbf{X} \text{diag}(\nu \circ \alpha)];$
 - 6 **Return** $\langle \nu, \mathbf{Z} \rangle;$
-

Robust solution of partitioned system $\mathbf{TX} = \mathbf{B}$.

Algorithm 4: $\langle \alpha, \mathbf{X} \rangle \leftarrow \text{RobustBlockedSolveMRHS}(\mathbf{T}, \mathbf{B})$

```
1 for  $(i, j) \in \{1, 2, \dots, M\} \times \{1, 2, \dots, N\}$  do
2    $\langle \alpha_{ij}, \mathbf{X}_{ij} \rangle \leftarrow \langle \mathbf{e}, \mathbf{B}_{ij} \rangle$ ;
3 for  $j \leftarrow M, M - 1, \dots, 1$  do
4   for  $k \leftarrow 1, 2, \dots, N$  do
5      $\langle \alpha_{jk}, \mathbf{X}_{jk} \rangle \leftarrow \text{RobustTRSM}(\mathbf{T}_{jj}, \langle \alpha_{jk}, \mathbf{X}_{jk} \rangle)$ ;
6     for  $i \leftarrow 1, 2, \dots, j - 1$  do
7       for  $k \leftarrow 1, 2, \dots, N$  do
8          $\langle \alpha_{ik}, \mathbf{X}_{ik} \rangle \leftarrow \text{RobustGEMM}(\langle \alpha_{ik}, \mathbf{X}_{ik} \rangle, \mathbf{T}_{ij}, \langle \alpha_{jk}, \mathbf{X}_{jk} \rangle)$ ;
9 for  $k \leftarrow 1, 2, \dots, N$  do
10    $\alpha_k \leftarrow \min_i \{ \alpha_{ik} \}$ ;
11   for  $i \leftarrow 1, 2, \dots, M$  do
12      $\mathbf{X}_{ik} \leftarrow \mathbf{X}_{ik} \text{diag}(\alpha_k \oslash \alpha_{ik})$ ;
13  $\alpha \leftarrow [\alpha_1 \quad \alpha_2 \quad \dots \quad \alpha_N]$ ;
14 Return  $\langle \alpha, \mathbf{X} \rangle$ ;
```

Performance degradation of robust solvers?

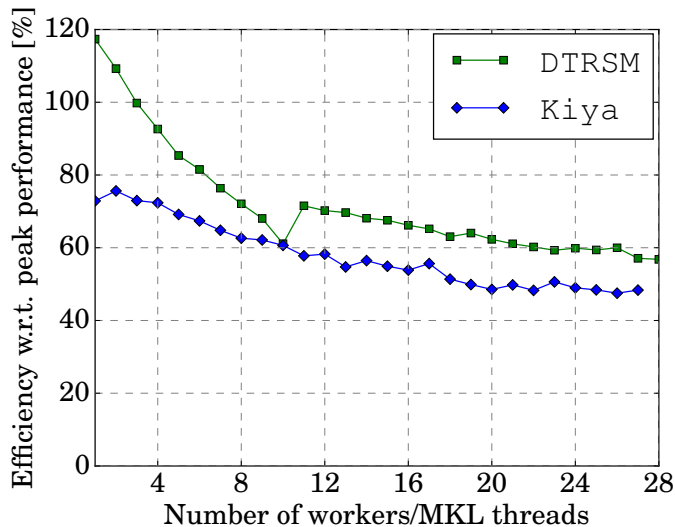


Figure: Triangular linear system $m = 40000$ and $n = 2000$

Caution: Underflow can still destroy the result!

$$\left[\begin{array}{c|cccc} \frac{1}{2} & & & & \\ \hline & M & & & \\ & -2 & 1 & & \\ & & \ddots & \ddots & \\ & & & -2 & 1 \end{array} \right] \begin{bmatrix} 2\Omega \\ \frac{\epsilon}{M} \\ \frac{2^1\epsilon}{M} \\ \vdots \\ \frac{2^{m-1}\epsilon}{M} \end{bmatrix} = \begin{bmatrix} \Omega \\ \epsilon \\ 0 \\ \vdots \\ 0 \end{bmatrix} .$$

New robust solvers

During the NLAFFET project, UMU has made robust solvers for

1. Standard linear systems, SM, StarPU:
2. Eigenvectors
 - 2.1 Standard case, real arithmetic, SM: StarNeig
 - 2.2 Generalized case, real arithmetic, SM: StarNeig
3. Sylvester matrix equations, SM, OpenMP.

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Lars Karlsson

Bo Kågström

Björn Adlerborn

Mahmoud Eljammally

Interesting problems which can now be solved ...

1. Catastrophic underflow in triangular equations can be avoided.
2. Proof of correctness of `ztgevc` which solves

$$\mathbf{SZ} = \mathbf{TZ}\mathbf{\Lambda}.$$

3. Proof of correctness of `dladiv` which does complex division.

$$p + iq = \frac{a + ib}{c + id}.$$