

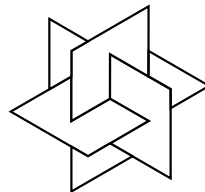
**Blocked and Multishift Variants
of the QZ Algorithm
for Computing Deflating Subspaces
of Regular Matrix Pencils**

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DFG research center Berlin
mathematics for key technologies

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Introduction

Want to solve generalized eigenvalue problem for **matrix pencil**

$$A - \lambda B, \quad A, B \in \mathbb{R}^{n \times n}.$$

This consists of:

- Finding **generalized eigenvalues** λ :

$$\det(A - \lambda B) = 0.$$

- Finding right and left **deflating subspaces** \mathcal{X} and \mathcal{Y} :

$$A\mathcal{X} \subseteq \mathcal{Y}, \quad B\mathcal{X} \subseteq \mathcal{Y}.$$



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Assumption: $A - \lambda B$ is a regular pencil, i.e., $\det(A - \lambda B) \not\equiv 0$.

The Basic QZ Algorithm

Moler/Stewart '73: QZ generates a sequence of orthogonally equivalent matrix pencils:

$$(A_0 - \lambda B_0) := (A - \lambda B), (A_1 - \lambda B_1), (A_2 - \lambda B_2), \dots$$

Under suitable conditions (Watkins/Elsner '94) :

$$(A_i - \lambda B_i) \longrightarrow \left(\begin{array}{c|c} \triangle & \triangle \\ \hline & \end{array} - \lambda \begin{array}{c|c} \triangle & \triangle \\ \hline & \end{array} \right).$$



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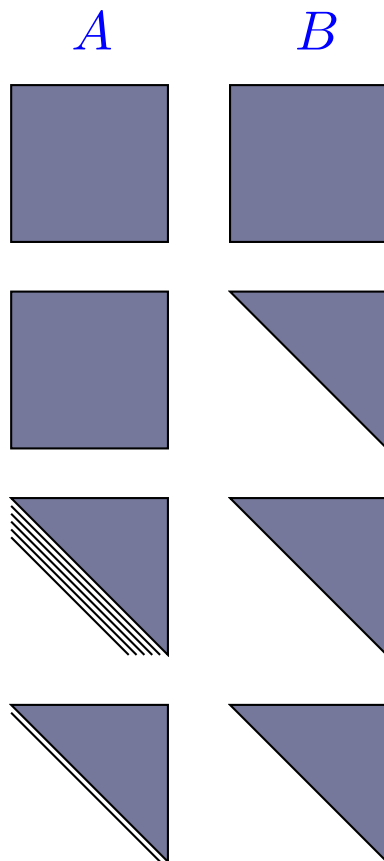
$$(A_i - \lambda B_i) \longrightarrow \left(\begin{array}{c|c} \triangle & \\ \hline & \triangle \end{array} - \lambda \begin{array}{c|c} \triangle & \\ \hline & \triangle \end{array} \right).$$

Three ingredients of implicit QZ:

- initial reduction to Hessenberg-triangular form;
- deflation;
- QZ iterations = bulge chasing.



Hessenberg-Triangular Reduction



original matrix pencil $A - \lambda B$

blocked QR factorization of B and update of A

blocked reduction to block Hessenberg-triangular form

reduction to Hessenberg-triangular form
based on pipelined Givens rotations

(Dackland/Kågström '99)



Up to **three times faster** than LAPACK's DGGHRD.



Deflation I: Small Subdiagonal Entry in A

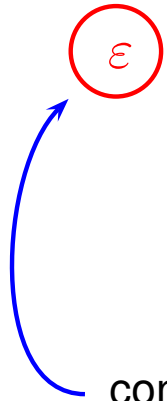
$$\begin{bmatrix} a & a & a & a & a & a \\ a & a & a & a & a & a \\ & a & a & a & a & a \\ & & \varepsilon & a & a & a \\ & & & a & a & a \\ & & & & a & a \end{bmatrix} - \lambda \begin{bmatrix} b & b & b & b & b & b \\ & b & b & b & b & b \\ & & b & b & b & b \\ & & & b & b & b \\ & & & & b & b \\ & & & & & b \end{bmatrix}$$





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considerably small if $|a_{j+1,j}| \leq \mathbf{u} \cdot (|a_{jj}| + a_{j+1,j+1}|)$





Deflation I: Small Subdiagonal Entry in A

$$\left[\begin{array}{ccc|ccc} a & a & a & a & a & a \\ a & a & a & a & a & a \\ & a & a & a & a & a \\ \hline & & 0 & a & a & a \\ & & & a & a & a \\ & & & & a & a \end{array} \right] - \lambda \left[\begin{array}{ccc|ccc} b & b & b & b & b & b \\ & b & b & b & b & b \\ & & b & b & b & b \\ \hline & & & b & b & b \\ & & & & b & b \\ & & & & & b \end{array} \right]$$

\Rightarrow generalized eigenvalue problem deflated into two smaller ones.





Deflation II: Small Diagonal Entry in B

$$\begin{bmatrix} a & a & a & a & a & a \\ a & a & a & a & a & a \\ & a & a & a & a & a \\ & & a & a & a & a \\ & & & a & a & a \\ & & & & a & a \end{bmatrix} - \lambda \begin{bmatrix} b & b & b & b & b & b \\ & b & b & b & b & b \\ & & \varepsilon & b & b & b \\ & & & b & b & b \\ & & & & b & b \\ & & & & & b \end{bmatrix}$$



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considerably small if $|b_{jj}| \leq \mathbf{u} \cdot (|b_{j-1,j}| + |b_{j,j+1}|)$





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$$\begin{bmatrix} a & a & a & a & a & a \\ a & a & a & a & a & a \\ & a & a & a & a & a \\ & & a & a & a & a \\ & & & a & a & a \\ & & & & a & a \end{bmatrix} - \lambda \begin{bmatrix} b & b & b & b & b & b \\ & b & b & b & b & b \\ & & 0 & b & b & b \\ & & & b & b & b \\ & & & & b & b \\ & & & & & b \end{bmatrix}$$





Deflation II: Small Diagonal Entry in B

$$\begin{bmatrix} a & a & a & a & a & a \\ 0 & a & a & a & a & a \\ & a & a & a & a & a \\ & & a & a & a & a \\ & & & a & a & a \\ & & & & a & a \end{bmatrix} - \lambda \begin{bmatrix} 0 & b & b & b & b & b \\ & b & b & b & b & b \\ & & b & b & b & b \\ & & & b & b & b \\ & & & & b & b \\ & & & & & b \end{bmatrix}$$

.. by a sequence of Givens rotations ..



Deflation II: Small Diagonal Entry in B

$$\left[\begin{array}{c|cccccc} a & a & a & a & a & a \\ \hline 0 & a & a & a & a & a \\ & a & a & a & a & a \\ & & a & a & a & a \\ & & & a & a & a \\ & & & & a & a \end{array} \right] - \lambda \left[\begin{array}{c|ccccc} 0 & b & b & b & b & b \\ \hline & b & b & b & b & b \\ & & b & b & b & b \\ & & & b & b & b \\ & & & & b & b \\ & & & & & b \end{array} \right]$$

\Rightarrow one eigenvalue $\lambda = \infty$ deflated.



Implicit QZ Iteration

Assumption: B is nonsingular.

Goal: Drive subdiagonal entries of A to ε .

Shift polynomial: Define shifts σ_1, σ_2 as generalized eigenvalues of the bottom right 2×2 subpencil of $A - \lambda B$. Let

$$x = (AB^{-1} - \sigma_1 I_n)(AB^{-1} - \sigma_2 I_n)e_1,$$

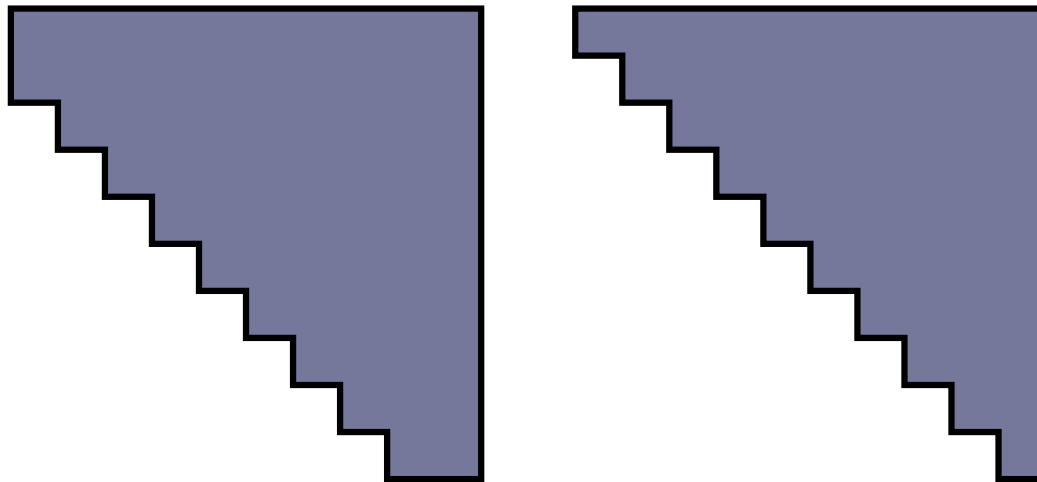
and Q such that $Q^T x = \alpha e_1$.

QZ iteration: Reduce $Q^T A - \lambda Q^T B$ back to Hessenberg-triangular form.

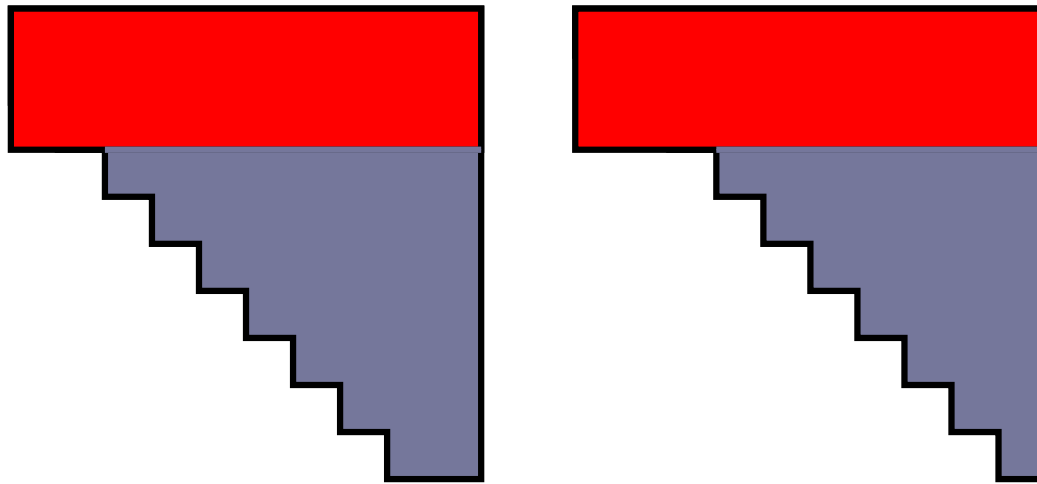




Bulge Chasing



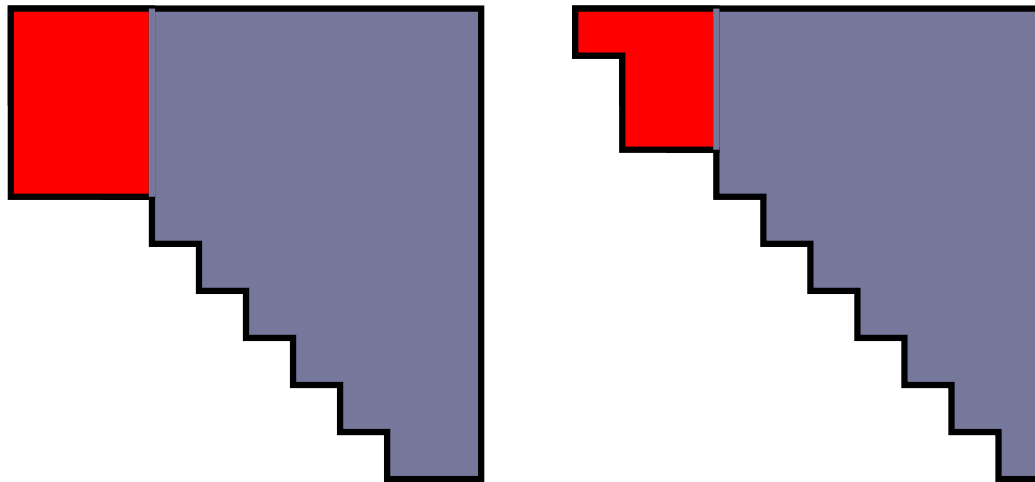
Bulge Chasing



..apply Q^T from the left..



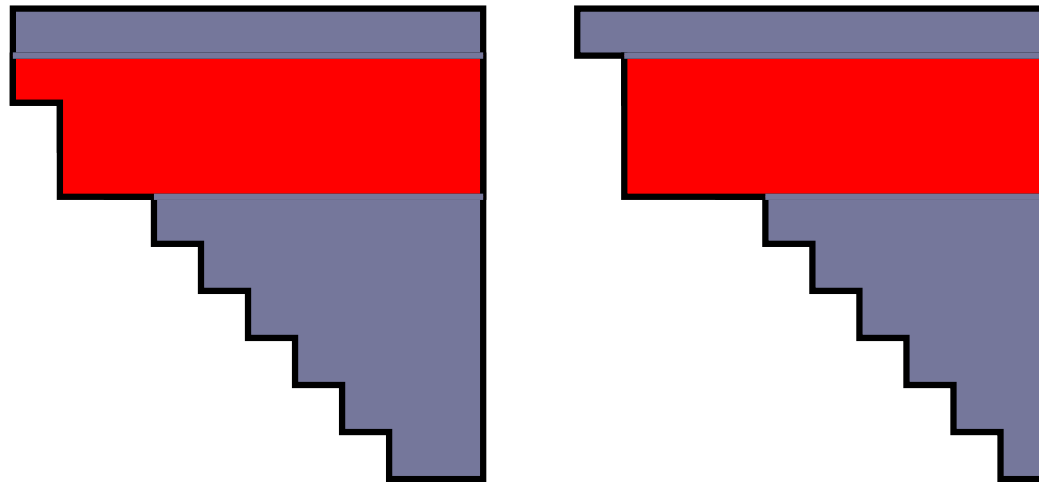
Bulge Chasing



..reduce 1st column of B by (opposite) Householder from the right..



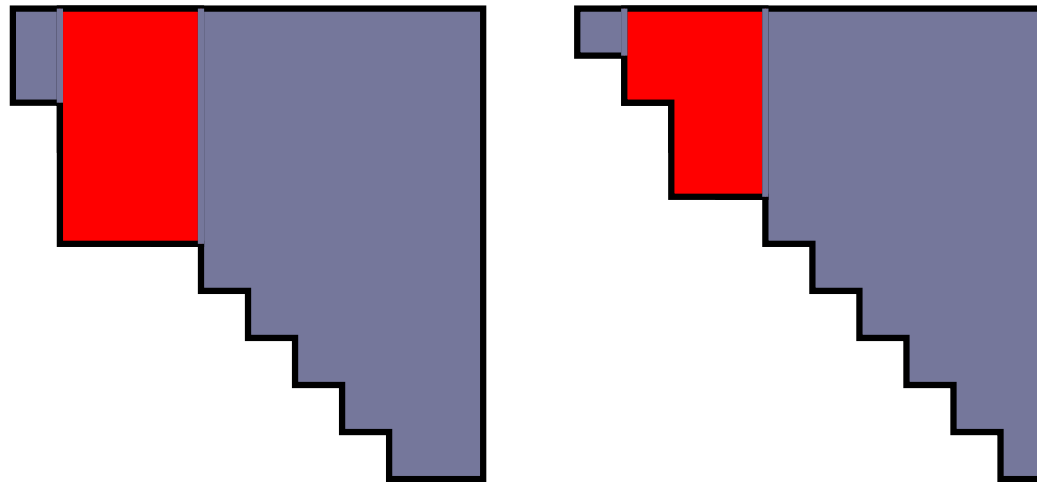
Bulge Chasing



..reduce 1st column of A by Householder from the left..



Bulge Chasing

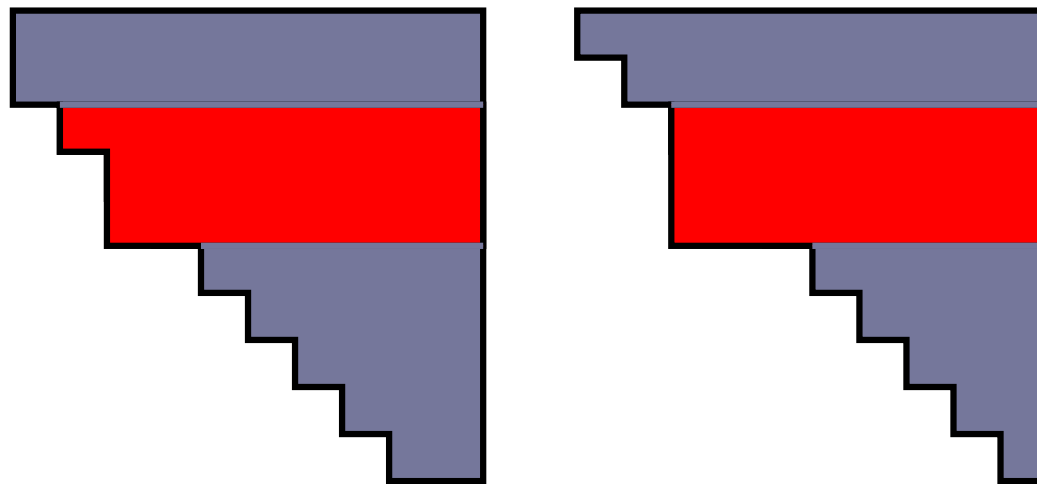


..reduce 2nd column of B by (opposite) Householder from the right..



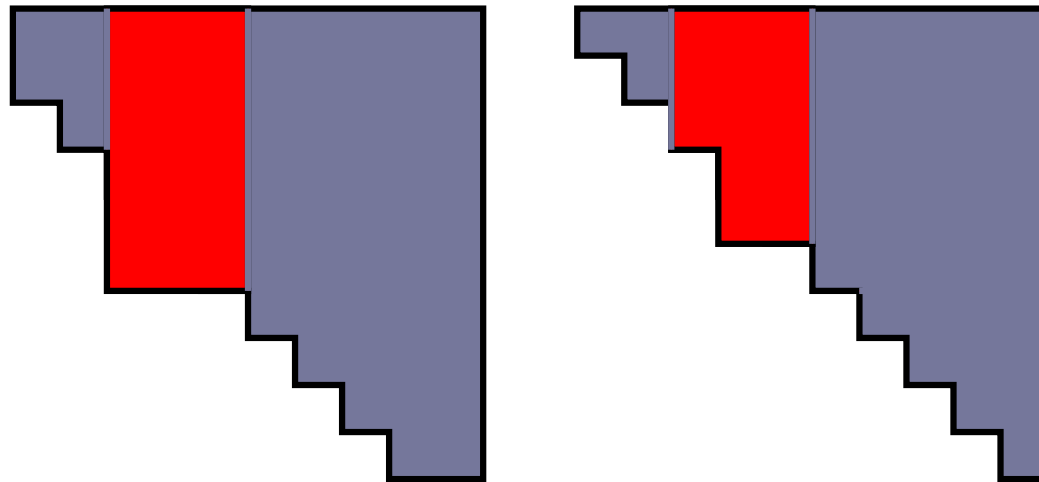


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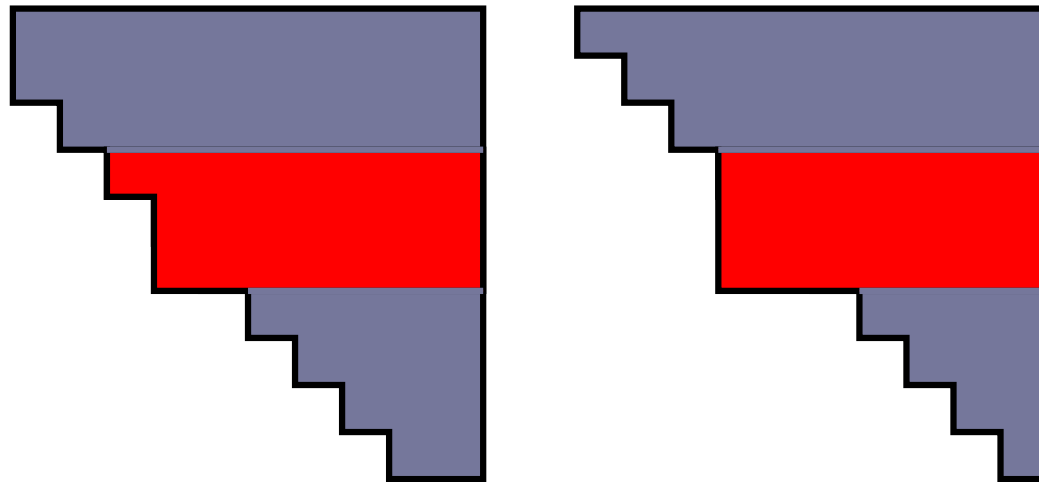


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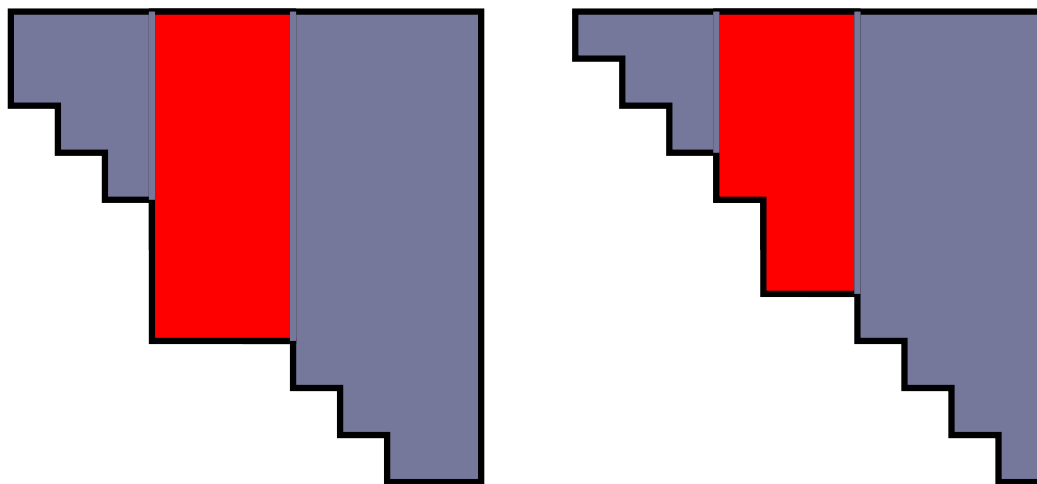


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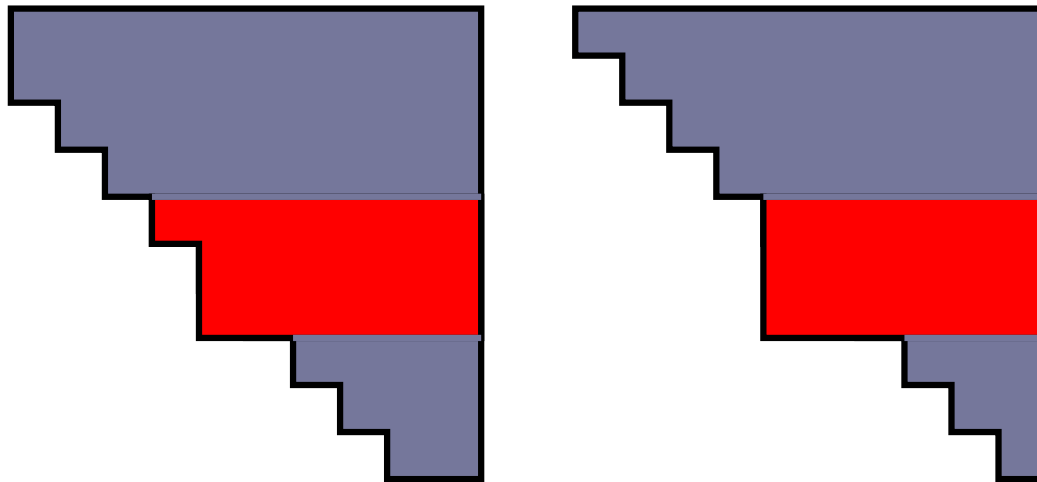




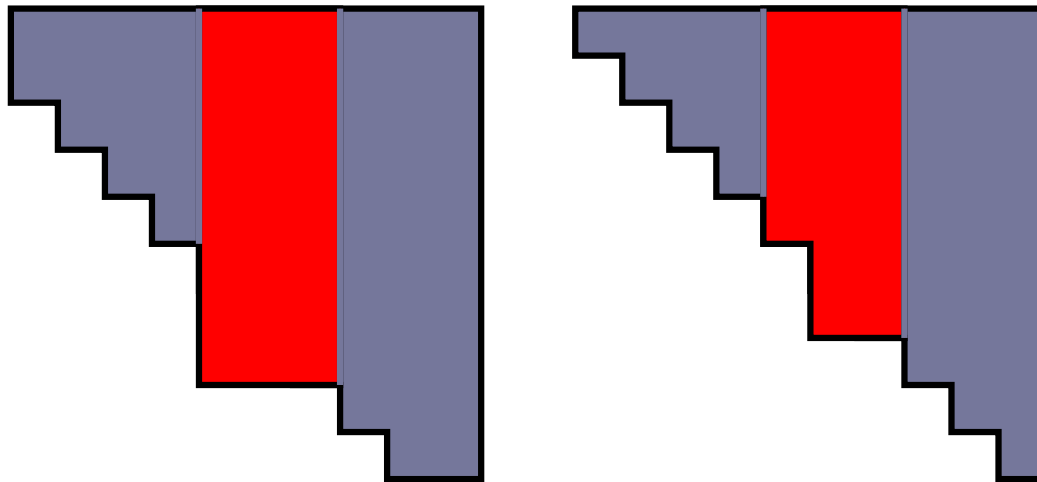
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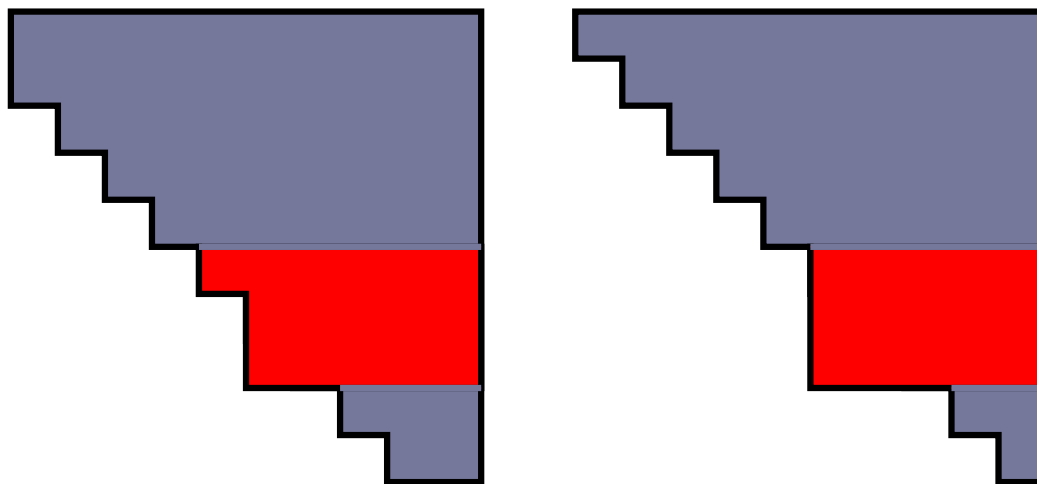


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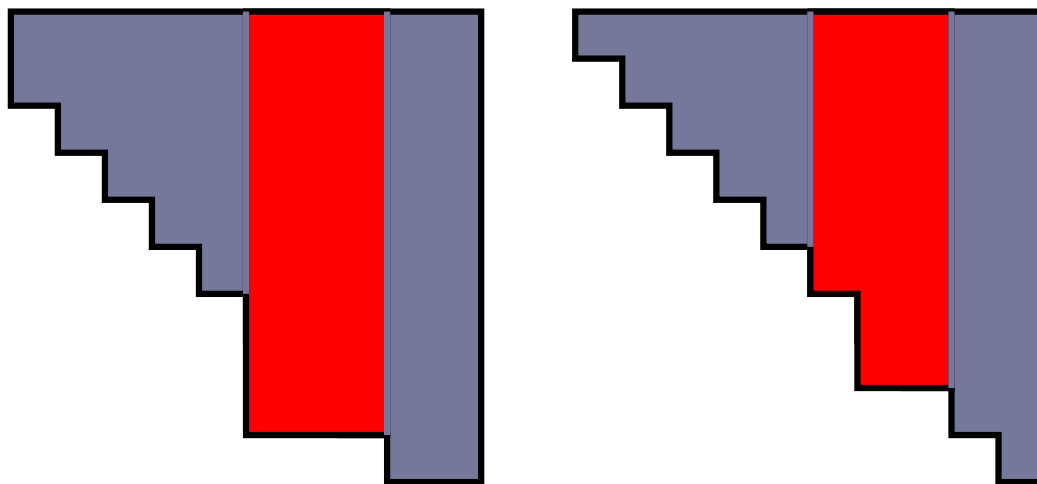


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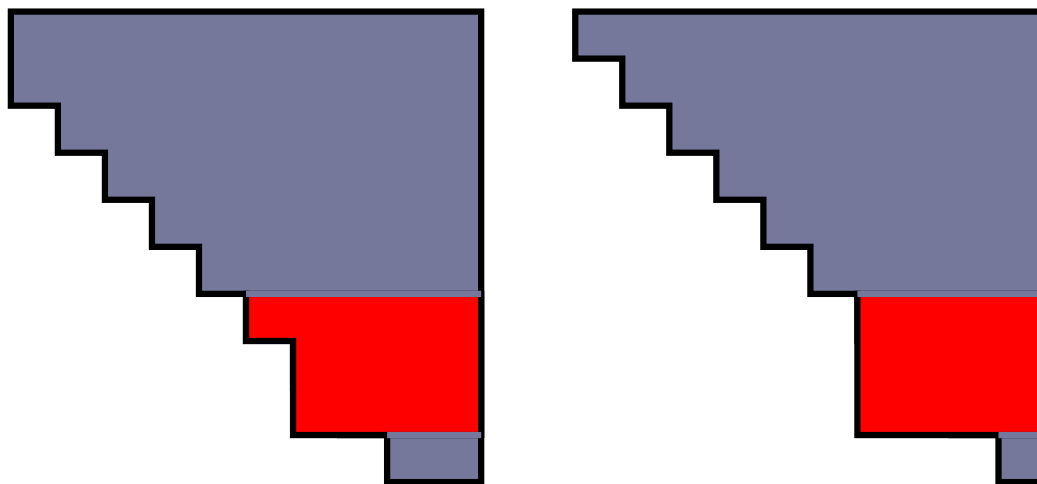


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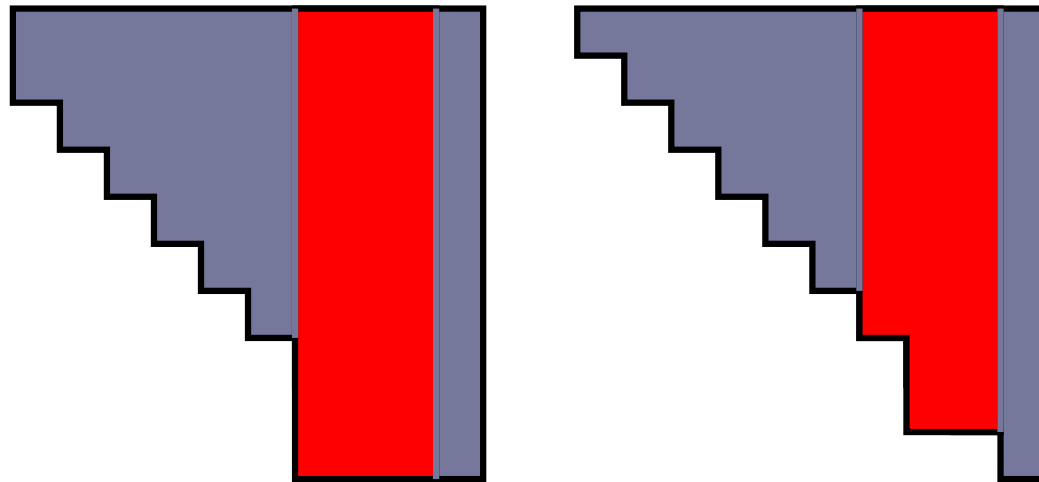


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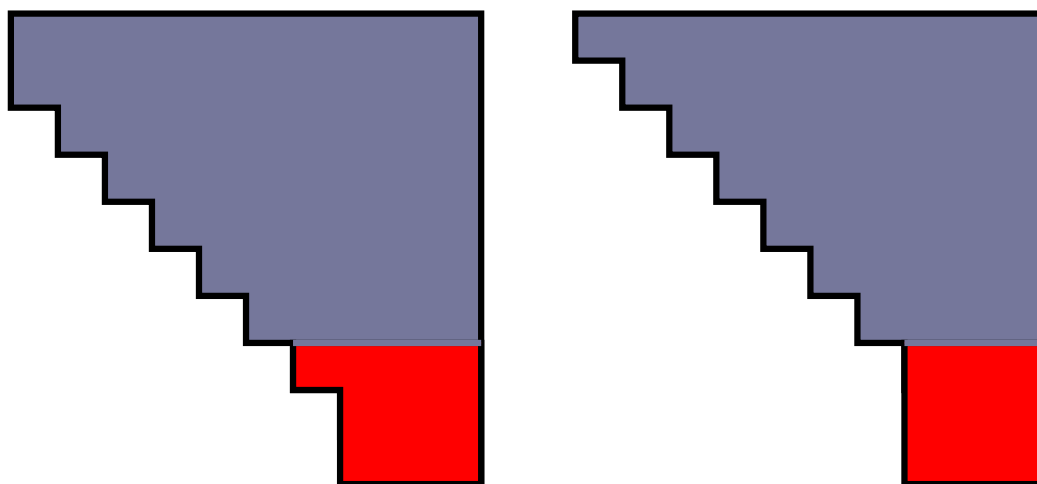


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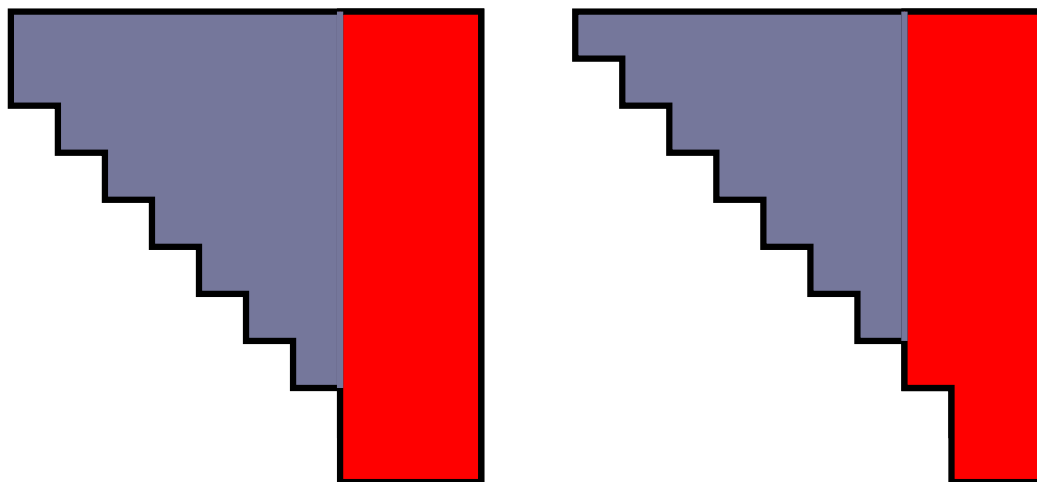


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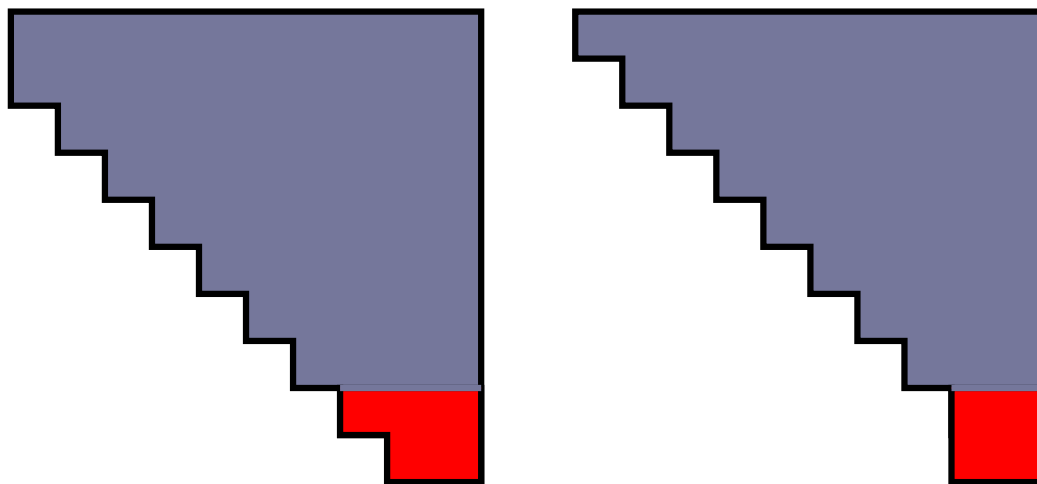


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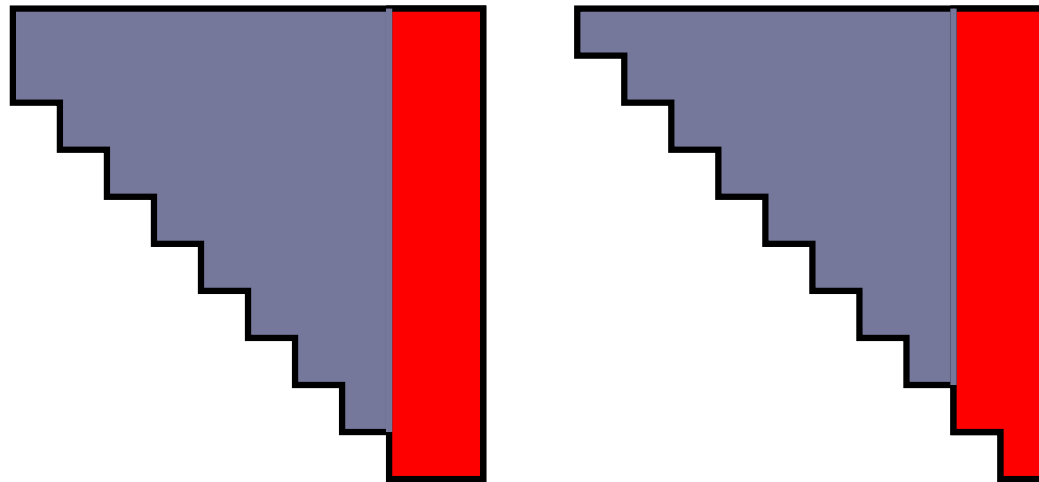


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- One QZ iteration requires $\mathcal{O}(n^2)$ flops.
- On average, roughly two QZ iterations are necessary for deflating a gen. eigenvalue (typically at the bottom right corner).
- High memory access/computation ratio and poor memory access pattern \Rightarrow poor performance!



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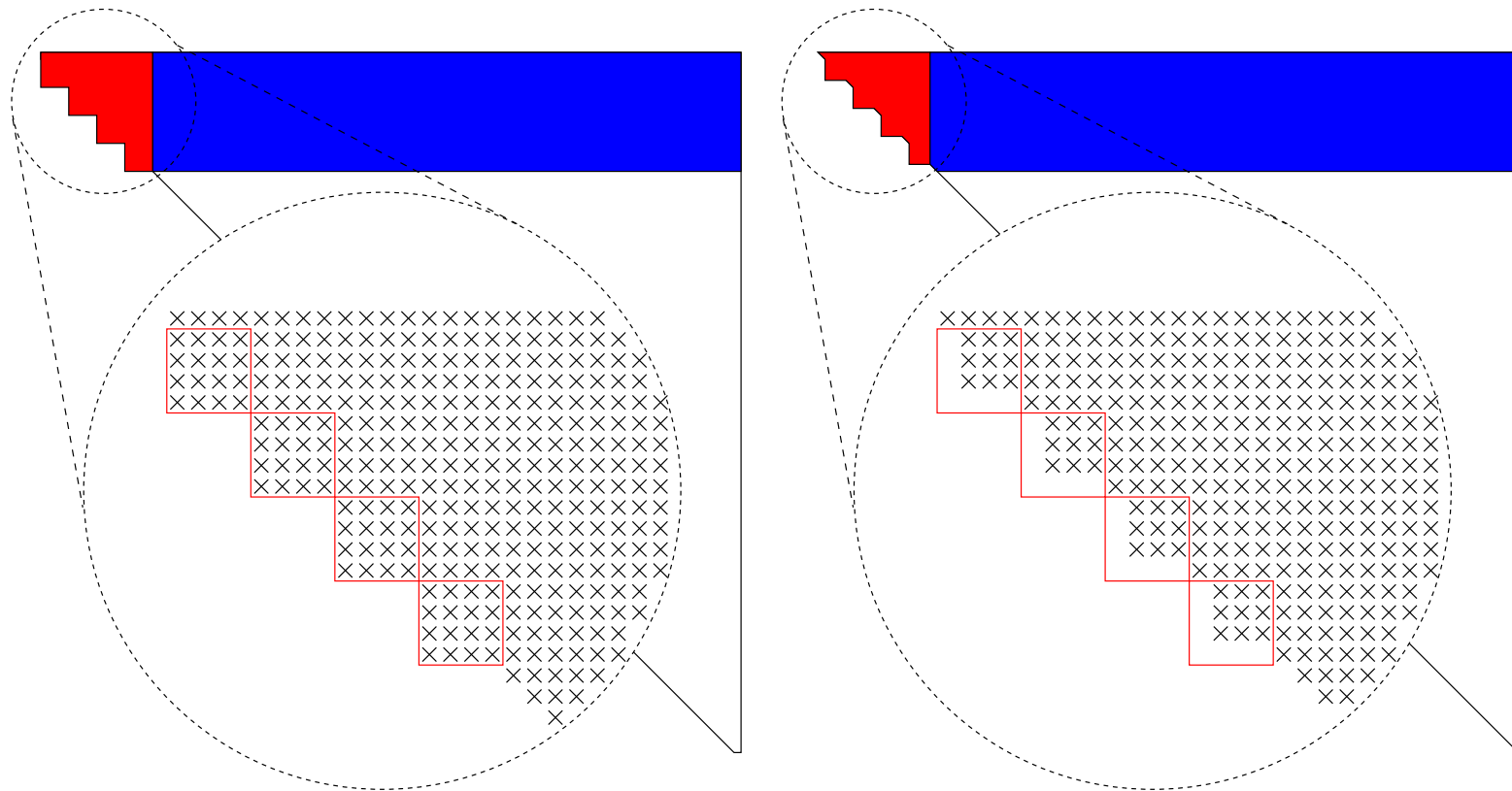
Remedy: Use more shifts per iteration.

But: Large bulge sizes lead to shift blurring phenomena and loss of convergence (Watkins '96, Kressner '04).

Use **tightly coupled chain of small bulges** instead. Based on ideas of Braman/Byers/Mathias '02, Lang '97, and many others for the QR algorithm.



Multishift QZ: Introducing a Chain of Bulges

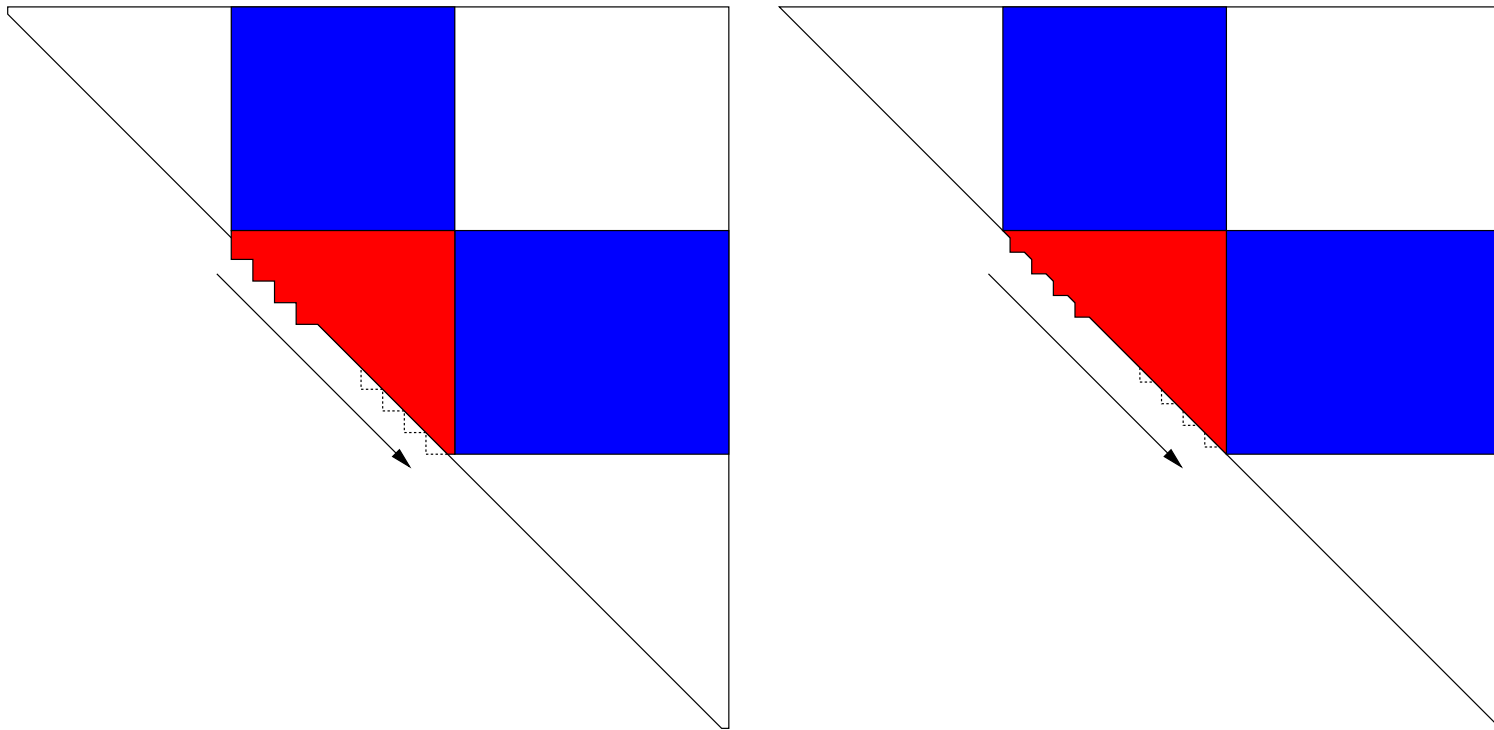


Red area: Updated during introduction.

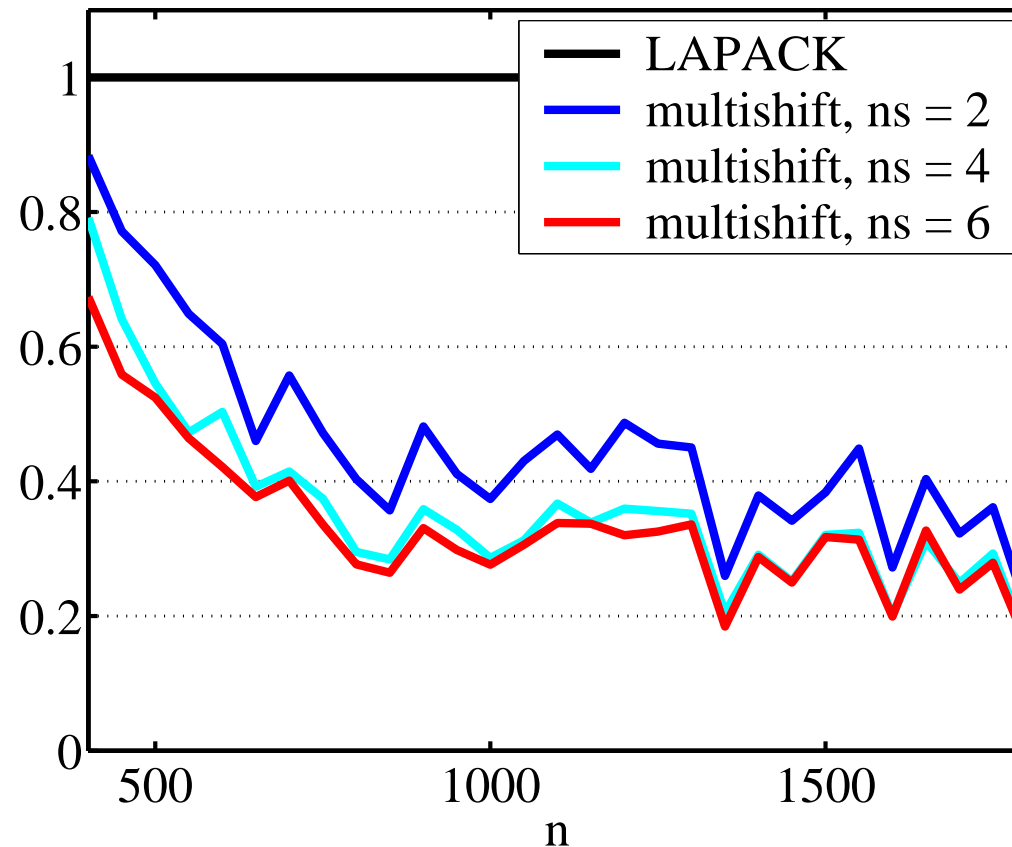
Blue area: Updated after introduction via matrix-matrix-mult.



Multishift QZ: Chasing a Chain of Bulges



Performance of the QZ algorithm

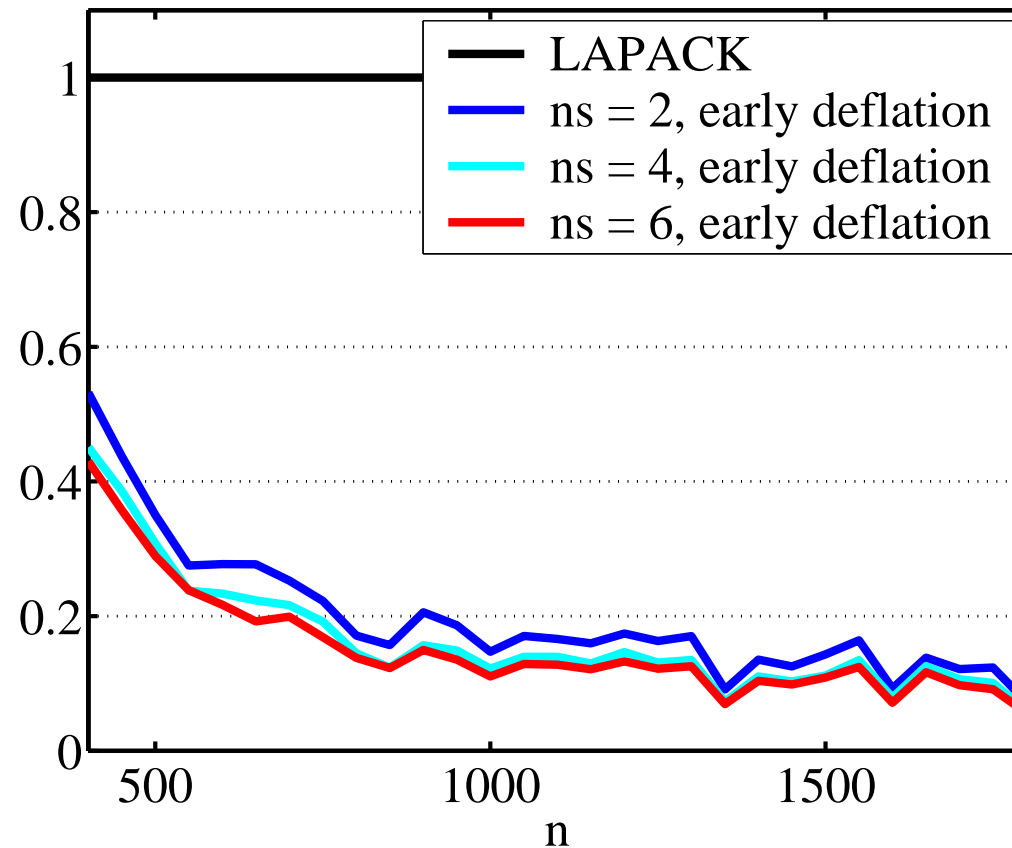


shifts/bulge: $n_s \in \{2, 4, 6\}$

shifts/QZ iteration: $m = 60$

Aggressive Early Deflation

Based on work by Braman/Byers/Mathias '02 for the QR algorithm, a more effective deflation strategy can be used to accelerate convergence of the QZ algorithm.



Computation of Deflating Subspaces

Output of QZ algorithm: Orthogonal matrices Q, Z such that

$$Q^T AZ - \lambda Q^T BZ = \begin{array}{c} \triangle \\ \triangle \end{array} - \lambda \begin{array}{c} \triangle \\ \triangle \end{array}.$$

First k columns of Q and Z span pair of deflating subspaces belonging to gen. eigenvalues of $k \times k$ leading principal subpencil.

To compute other deflating subspaces, gen. eigenvalues must be reordered.

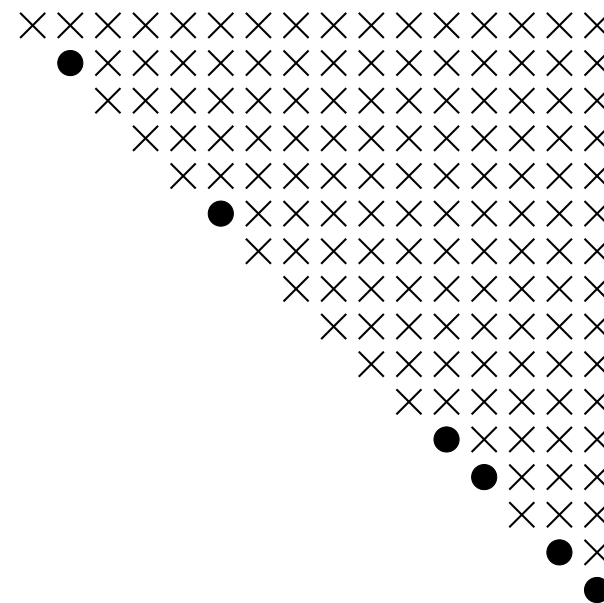
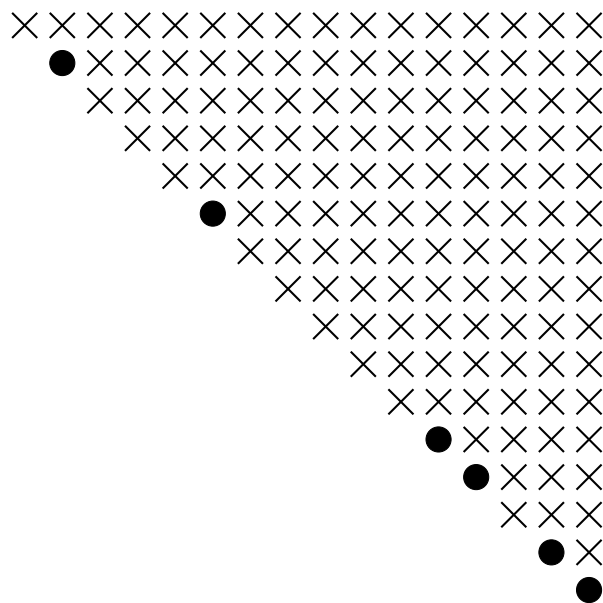
Van Dooren '82, Kågström '93, Kågström/Poromaa '96, propose to reorder gen. eigenvalues in a bubble sort-like fashion.

Again: High memory access/computation ratio and poor memory

access pattern \Rightarrow poor performance!

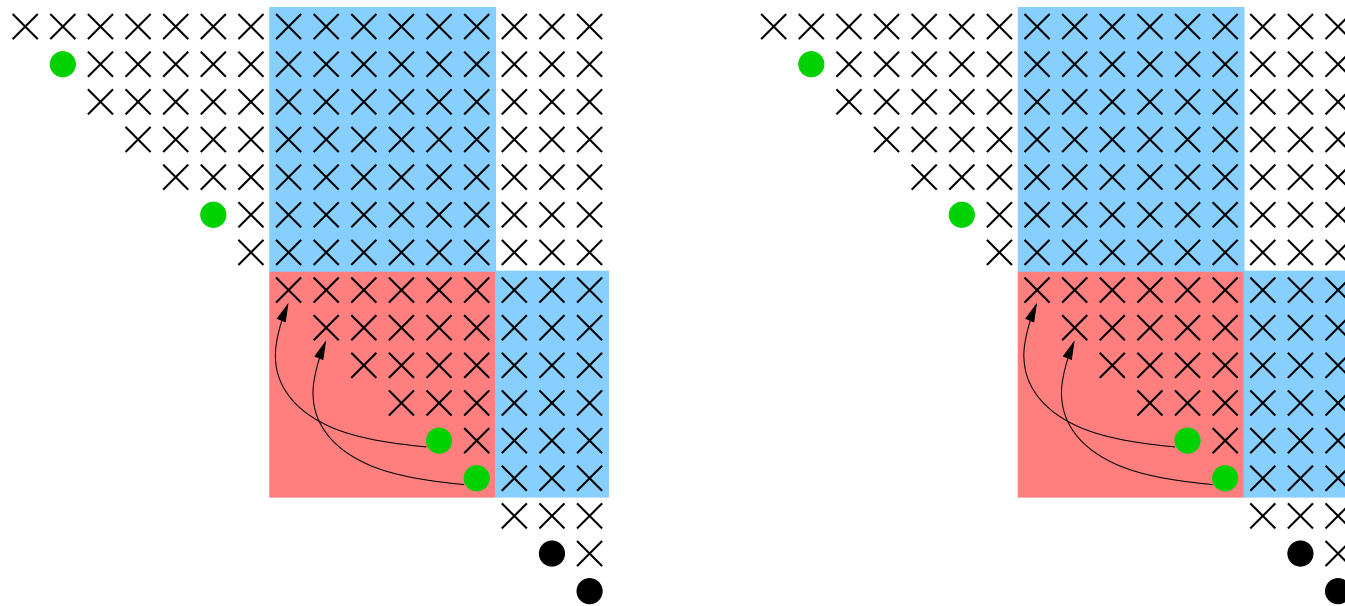


Block Algorithm for Reordering Eigenvalues



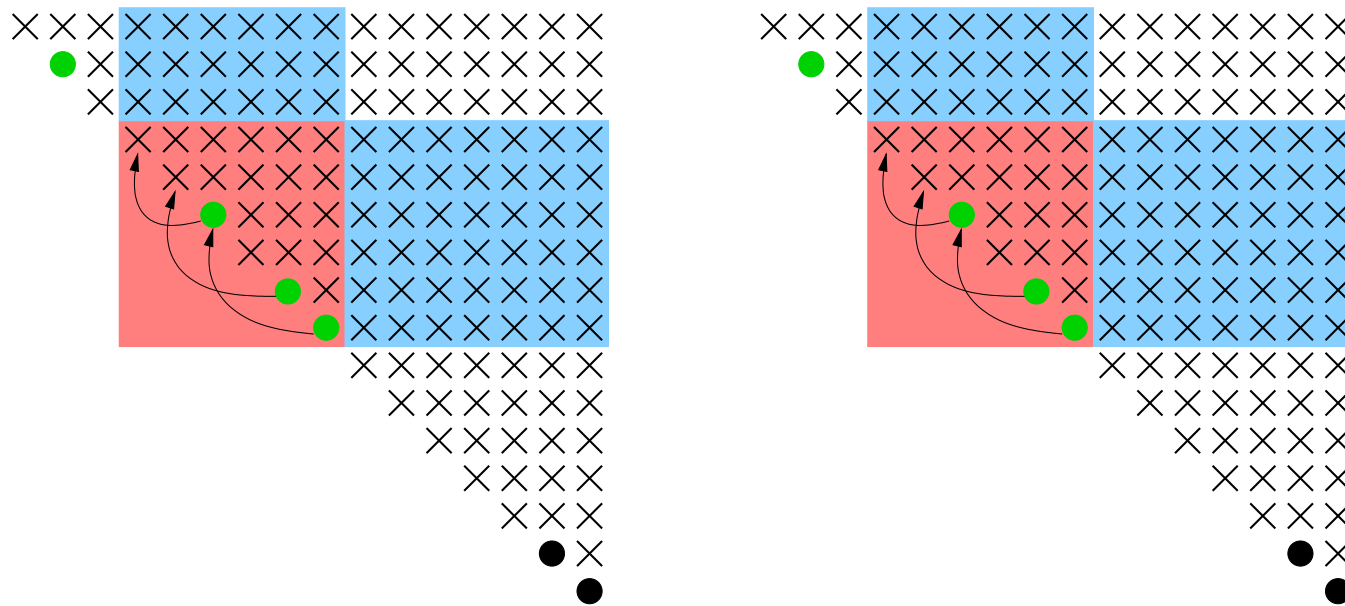


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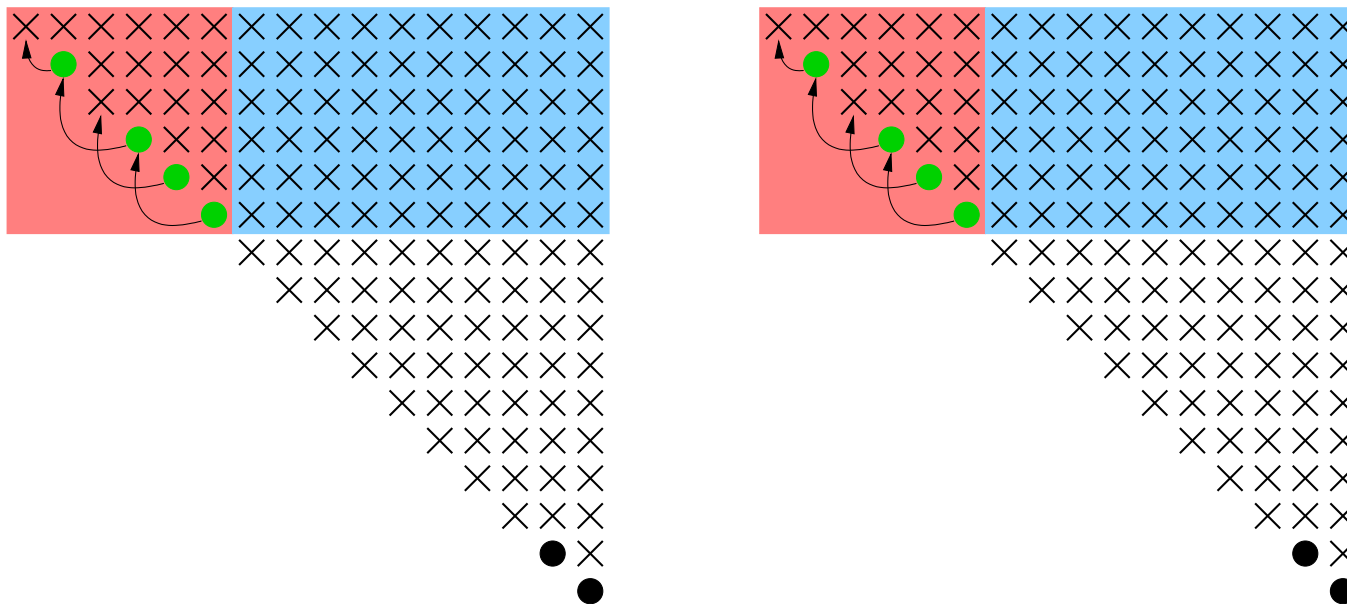


Block Algorithm for Reordering Eigenvalues





Block Algorithm for Reordering Eigenvalues





Performance of Block Reordering Algorithm

For **standard** eigenvalue problem:

n	sel.	LAPACK	new	ratio
500	5%	0.25	0.09	36%
500	25%	0.75	0.25	33%
500	50%	0.81	0.33	40%
1000	5%	2.87	0.60	21%
1000	25%	8.40	1.57	19%
1000	50%	10.08	2.10	21%
1500	5%	9.46	1.69	18%
1500	25%	30.53	4.88	16%
1500	50%	35.93	6.55	18%



Concluding Remarks



Work under progress:

- Integration of described algorithms into new release of LAPACK.
- ScaLAPACK-like parallel implementation of QZ algorithm (Björn Adlerborn, Univ. Umeå)



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For more details, see:

- Adlerborn/Dackland/Kågström: Parallel and blocked algorithms for reduction of a regular matrix pair to Hessenberg-triangular and generalized Schur forms. PARA2002, Springer-Verlag, LNCS, Vol. 2367, pp 319–328.
- Kressner/Kågström: Multishift variants of the QZ algorithm with aggressive early deflation. In preparation, 2004.
- Kressner: Numerical algorithms and software for general and structured eigenvalue problems. PhD thesis, 2004.

